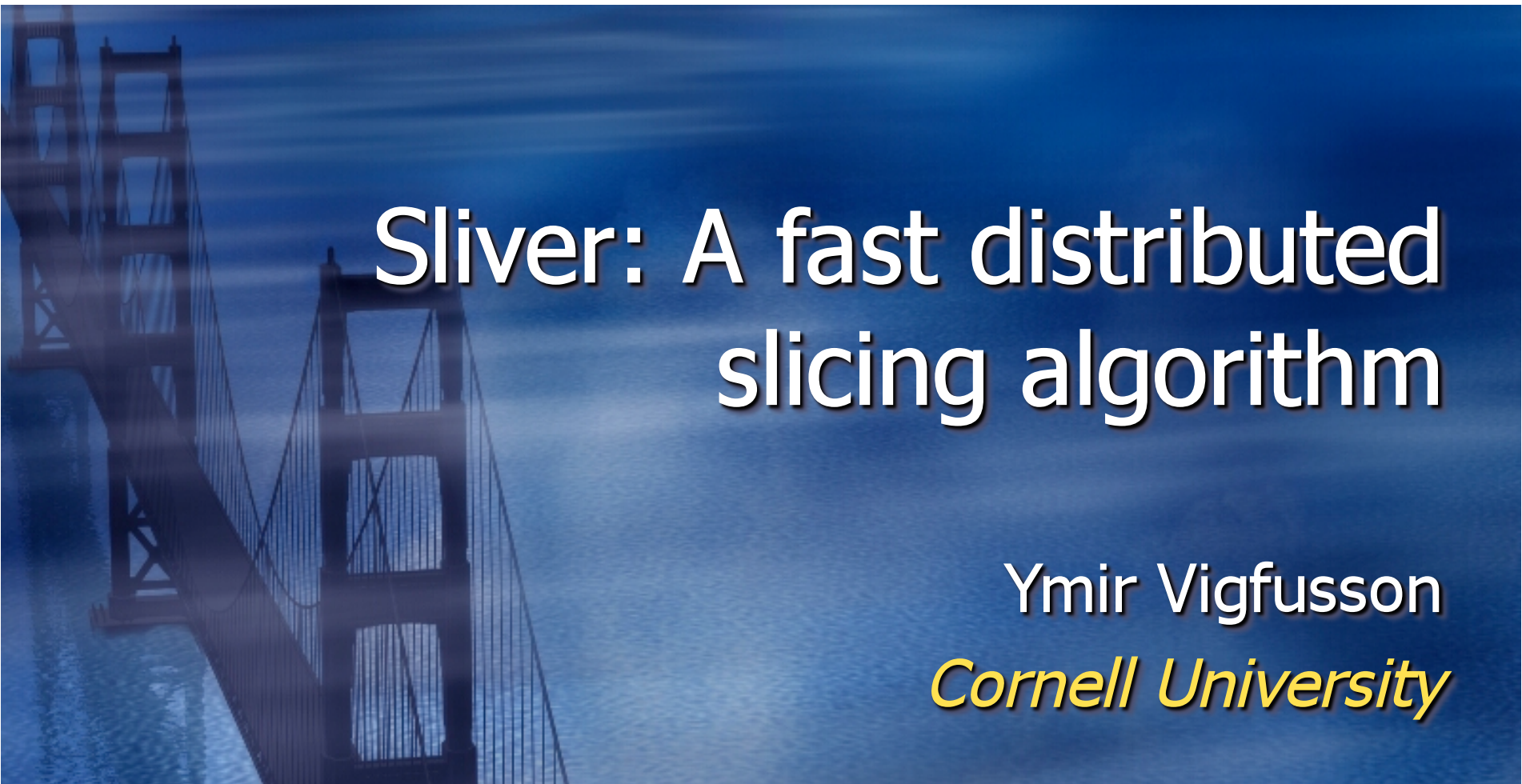


A background image of the Golden Gate Bridge in San Francisco, viewed from a low angle looking up at the suspension towers and cables. The bridge is silhouetted against a blue sky with light clouds.

Sliver: A fast distributed slicing algorithm

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The Distributed Slicing Problem

- ♦ n nodes, each has an attribute value x_i

1 4 5 7 50 80 83 95 98

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I'm in slice 2!



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- ◆ Node j estimates its slice as $\left\lceil \frac{k B_j}{m} \right\rceil$

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- ♦ Experiments on Emulab and simulations on Skype traces indicate rapid convergence and churn-tolerance
- ♦ *Conclusion:* Sliver is simple and robust with fast convergence properties.